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DISTRIBUTION-FREE PERFORMANCE BOUNDS WITH THE RESUBSTITUTION ER--ETC(U)

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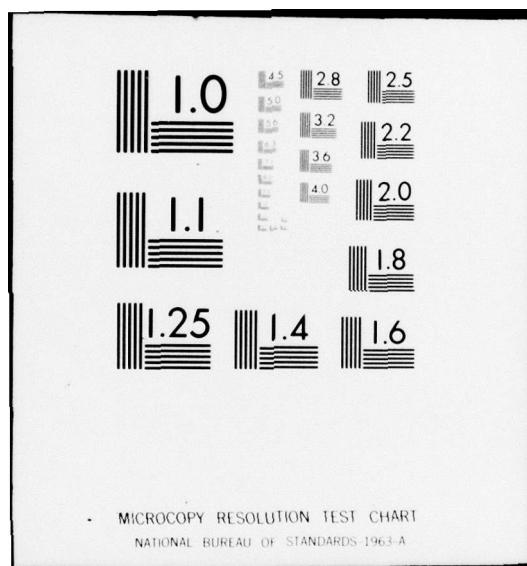
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DISTRIBUTION-FREE PERFORMANCE BOUNDS WITH
THE RESUBSTITUTION ERROR ESTIMATELuc P. Devroye and T.J. Wagner*
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Probability inequalities are given for the deviation of the resubstitution error estimate from the unknown conditional probability of error. The inequalities are distribution-free and can be applied to linear discrimination rules, to nearest neighbor rules with a reduced sample size, and to histogram rules.

1. Introduction

The discrimination problem may be formulated as follows. The statistician collects data $(X_1, \theta_1), \dots, (X_n, \theta_n)$, a sequence of independent identically distributed random vectors drawn from the distribution of (X, θ) , a random vector independent of the data. For each $1 \leq i \leq n$, the observation X_i takes values in \mathbb{R}^m and its state θ_i takes values in $\{1, \dots, M\}$. The discrimination problem is that of estimating the state θ from the data and the observation X using procedures which do not require complete knowledge of the distribution of (X, θ) . If $\hat{\theta}$ denotes the estimate, that is, $\hat{\theta} = g(X, V_n)$ where g is a Borel measurable $\{1, \dots, M\}$ -valued function of X and the data $V_n = (X_1, \theta_1), \dots, (X_n, \theta_n)$, then a measure of the performance of the procedure given the data is $L_n = P\{\hat{\theta} \neq \theta | V_n\}$, the conditional probability of error.

Since the distribution of (X, θ) is unknown, there is in general no way of computing L_n from the data. Using the data the statistician may try to estimate L_n by \hat{L}_n . A survey of estimation techniques can be found in Toussaint.¹ One of the oldest estimates is the resubstitution estimate

$$\hat{L}_n = n^{-1} \sum_{i=1}^n I\{\hat{\theta}_i \neq \theta_i\}$$

where $\hat{\theta}_i = g(X_i, V_n)$, $1 \leq i \leq n$, are the estimates of the states of X_1, \dots, X_n with the given discrimination procedure, and where I is the indicator function.

In this paper we obtain upper-bounds for $P\{|\hat{L}_n - L_n| \geq \epsilon\}$ that do not depend upon the distribution of (X, θ) , and that are applicable to three large classes of discrimination rules,

- (i) the linear discrimination rules,
- (ii) the nearest-neighbor rules with reduced sample size, and
- (iii) the histogram decision rules.

The existence of distribution-free bounds with the resubstitution estimate for linear discrimination rules was first noticed by Vapnik and Chervonenkis.¹⁰ The bounds for the class (i) improve the bounds given in Devroye and Wagner², while the results for the rules (ii) and (iii) are new. The possible existence of distribution-free bounds for (ii) was suggested to the authors by Dr. Penrod¹⁶.

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2. Main Results

Let $\phi_0, \phi_1, \dots, \phi_m$ be known measurable mappings from \mathbb{R}^m to \mathbb{R} where $m' \geq 1$ and $m \geq 1$, and $\phi_0 \equiv 1$. Let $w_0 = (w_{10}, \dots, w_{1m}), \dots, w_M = (w_{M0}, \dots, w_{Mm})$ be Borel-measurable $\mathbb{R}^{m'+1}$ -valued vector functions of the data V_n . Then, the rule which assigns the state $\hat{\theta} = j$ ($1 \leq j \leq M$) to X whenever j is the first integer for which

$$\sum_{i=0}^{m'} w_{ji}(V_n) \phi_i(X) = \max_{1 \leq k \leq M} \left\{ \sum_{i=0}^{m'} w_{ki}(V_n) \phi_i(X) \right\}$$

is called a linear discrimination rule (see Duda and Hart³ for a survey of the literature on linear discrimination). We emphasize that the w_1, \dots, w_M may be picked in an arbitrary fashion, using any method that can or cannot be found in the literature. The functions ϕ_i are picked in advance. The following bound is proved in the Appendix.

Theorem 1. For every $\epsilon > 0$ and for all linear discrimination rules with given $\phi_0, \phi_1, \dots, \phi_m$, the resubstitution estimate \hat{L}_n satisfies

$$P\{|\hat{L}_n - L_n| \geq \epsilon\} \leq 4M(1+(2n)^{m'})^{M-1} e^{-n\epsilon^2/8M^2}.$$

For the interesting case that $M=2$, we see that

$$P\{|\hat{L}_n - L_n| \geq \epsilon\} \leq 8(1+(2n)^{m'}) e^{-n\epsilon^2/32}.$$

Using the Borel-Cantelli lemma and Theorem 1, we see that for a given m' and M , and uniformly over all linear discrimination rules, $|\hat{L}_n - L_n| \not\equiv 0$ with probability one, a result due to Glick¹¹. Thus, the statistician could pick the w_1, \dots, w_M that minimize the resubstitution estimate \hat{L}_n because he knows from Theorem 1 that the corresponding probability of error L_n will be very close to \hat{L}_n , and that for large n , minimizing \hat{L}_n is nearly equivalent to minimizing L_n (see Wagner¹²).

In the literature special attention has been given to the nearest-neighbor rule with a reduced number of observations where the reduction is a result of editing (see for instance Wilson⁴), condensing (Hart⁵) or any other operation (Tomek⁶). In general, we end up with $\mathbb{R}^m \times \{1, \dots, M\}$ -valued random vectors $(Y_1, \xi_1), \dots, (Y_K, \xi_K)$ where K is an integer-valued random variable with $1 \leq K \leq k_n$. The (Y_j, ξ_j) and K may depend upon the data in an arbitrary fashion. A new observation X is assigned the state $\hat{\theta} = \xi_j$ whenever j is the smallest index for which

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$$\|X - Y_j\| = \min_{1 \leq i \leq K} \|X - Y_i\|.$$

Thus, $\hat{\theta}$ is the state of the nearest neighbor to X among Y_1, \dots, Y_K . In the Appendix the following Theorem is proved.

Theorem 2. For every $\epsilon > 0$, and for all the nearest-neighbor rules with reduced sample size, the resubstitution estimate \hat{L}_n satisfies

$$P(|\hat{L}_n - \hat{L}_n| \geq \epsilon) \leq 4M(1+(2n)^m)^{k_n-1} e^{-ne^2/8k_n^2}$$

where k_n is an upper-bound on the reduced sample size.

We remark that this bound is independent of the distribution of (X, θ) . The bound converges to 0 as n grows large provided that the sequence k_1, k_2, \dots is picked in such a way that $k_n^3 \log n/n \not\rightarrow 0$. It is clear that this bound is useless for the well-known nearest neighbor rule⁷, that is, the rule with $K=k_n=n$ and $(Y_i, \varepsilon_i) = (X_i, \theta_i)$, $1 \leq i \leq n$. This was to be expected because the resubstitution estimate with the nearest-neighbor rule is overly optimistic. In fact, if the probability measure μ of X is absolutely continuous with respect to Lebesgue measure, then $\hat{L}_n = 0$ with probability one, no matter what value \hat{L}_n takes.

Theorem 2 can be useful for reduced, selective, condensed or edited nearest-neighbor rules^{4-6,13-14}. If k_n is a prespecified number of (Y_i, ε_i) 's that are to be used in the new nearest-neighbor rule, then the statistician could compute \hat{L}_n with some selected data $(X_{i_1}, \theta_{i_1}), \dots, (X_{i_{k_n}}, \theta_{i_{k_n}})$ where $\{i_1, \dots, i_{k_n}\}$ is a sub-

set of $\{1, \dots, n\}$, and decide to use that set of indices for which the resubstitution estimate is minimal. Using Theorem 2, we also know how much confidence we can put in our estimate \hat{L}_n regardless of the selection procedure of the (Y_i, ε_i) and without any knowledge of the distribution of (X, θ) .

The (Y_i, ε_i) , $1 \leq i \leq k_n$, partition \mathbb{R}^m into k_n disjoint sets A_1, \dots, A_{k_n} where the state of X is estimated by $\hat{\theta} = \varepsilon_j$ whenever X takes values in A_j (that is, X is closest to Y_j). The partition in this case depends on the data because the Y_i depend upon the data. For a given fixed partition of \mathbb{R}^m , we can expect to obtain tighter upper-bounds for $P(|\hat{L}_n - \hat{L}_n| \geq \epsilon)$ even if the partition is not generated by a reduced nearest-neighbor rule.

Let A_1, \dots, A_{k_n} be any fixed partition of \mathbb{R}^m and let $\varepsilon_1, \dots, \varepsilon_{k_n}$ be $\{1, \dots, M\}$ -valued random variables where, as before, ε_j is the state assigned to X whenever X takes values in A_j . Such rules will be called histogram decision rules. We prove the following four distribution-free inequalities that are valid no matter how the ε_i depend upon the data. The inequalities

do not imply one another.

Theorem 3. For a given k_n -member partition of \mathbb{R}^m , for any way of specifying $\varepsilon_1, \dots, \varepsilon_{k_n}$ from the data in a histogram decision rule, and for every $\epsilon > 0$, the resubstitution estimate \hat{L}_n satisfies

$$P(|\hat{L}_n - \hat{L}_n| \geq \epsilon) \leq g_{n1}, \quad 1 \leq i \leq 4,$$

where

$$g_{n1} = 4 k_n \min(1+2n, M) e^{-ne^2/8k_n^2}$$

$$g_{n2} = 2 k_n M e^{-2ne^2/M^2k_n^2}$$

$$g_{n3} = 4 \min(M^k, 2^{2n}, (4n/k_n)^{k_n}) e^{-ne^2/8}$$

$$g_{n4} = 2 M e^{-2ne^2/M^2k_n^2}.$$

We note here that g_{n1} and g_{n3} are useful even if $M = \infty$ (i.e., the ε_i and θ_i can take a countably infinite number of values). Clearly, all the g_{ni} are independent of the dimension m and the distribution of (X, θ) , and $g_{n3} \not\rightarrow 0$ provided that $k_n/n \not\rightarrow 0$. If $k_n = \infty$, the bounds are not applicable, and, as we will see, the resubstitution estimate does not possess the distribution-free properties that it has with finite partitions of \mathbb{R}^m . Assume that A_1, A_2, \dots is a fixed countably infinite partition of \mathbb{R}^m . If the ε_i are random variables that are independent of the data, then

$$P(|\hat{L}_n - \hat{L}_n| \geq \epsilon) \leq 2e^{-2ne^2} \quad (1)$$

for any $\epsilon > 0$. However, such rules are impractical. The closest one can come to the Bayes rule with a fixed partition is to let $\varepsilon_i = j$ if j is the smallest integer such that $N_{ij} = \max_{1 \leq l \leq M} N_{il}$ where N_{ij} is the number of (X_k, θ_k) 's with $X_k \in A_i$ and $\theta_k = j$. But even with this obvious choice of $\varepsilon_1, \varepsilon_2, \dots$ we see that for any m and $M \geq 2$, there always exists a distribution of (X, θ) such that $|\hat{L}_n - \hat{L}_n| \geq \frac{1}{2}$ with probability one. Indeed, assume that $M=2$, that $\theta=2$ with probability one, and that X takes values in each A_1, \dots, A_{2n} with equal probability $1/2n$. If the ε_i are picked as described, then the resubstitution estimate \hat{L}_n equals 0. Furthermore,

$$\hat{L}_n = \sum_{i=1}^{2n} P(X \in A_i) I_{\{N_{i2}=0\}} \geq n/2n = 1/2.$$

This shows that even with the most obvious dependency of the ε_i on the data, we will never be able to upper-bound $P(|\hat{L}_n - \hat{L}_n| \geq \epsilon)$ by an expression that decreases to 0 as n grows large, uniformly over all distributions of (X, θ) .

3. Appendix

Proof of Theorem 1.

Let v be the probability measure of (X, θ) where X takes values in \mathbb{R}^m and θ takes values in $\{1, \dots, M\}$. It is clear that if v_n is the empirical measure for $(x_1, \theta_1), \dots, (x_n, \theta_n)$, and if A_1, \dots, A_M is the partition of \mathbb{R}^m that is generated by the linear discrimination rule (that is, A_i is the set on which we estimate the state of X by i), then

$$L_n = \sum_{i=1}^M v(A_i \times \{i\}^c)$$

and

$$\hat{L}_n = \sum_{i=1}^M v_n(A_i \times \{i\}^c).$$

Thus,

$$\begin{aligned} |L_n - \hat{L}_n| &= \left| \sum_{i=1}^M (v(A_i \times \{i\}^c) - v_n(A_i \times \{i\}^c)) \right| \\ &= \left| \sum_{i=1}^M (v_n(A_i \times \{i\}) - v_n(A_i \times \{i\})) \right| \\ &\leq M \sup_{\substack{A \in \mathcal{A} \\ i \in \{1, \dots, M\}}} |v_n(A \times \{i\}) - v(A \times \{i\})| \end{aligned}$$

where \mathcal{A} is the class of all sets that are intersections of $(M-1)$ linear halfspaces of \mathbb{R}^m . We recall that a linear halfspace of \mathbb{R}^m is a set of $x = (x^1, \dots, x^m)$ for which $x^1 a_1 + \dots + x^m a_m \geq a_0$ or $x^1 a_1 + \dots + x^m a_m < a_0$ for some $(a_0, a_1, \dots, a_m) \in \mathbb{R}^{m+1}$. Thus, every (a_0, a_1, \dots, a_m) defines two linear halfspaces.

By an inequality of Vapnik and Chervonenkis⁹,

$$\begin{aligned} P(|L_n - \hat{L}_n| \geq \epsilon) &\leq P(\sup_{\substack{A \in \mathcal{A} \\ 1 \leq i \leq M}} |v_n(A \times \{i\}) - v(A \times \{i\})| \geq \epsilon/M) \\ &\leq 4s(\mathcal{B}, 2n)e^{-n(\epsilon/M)^2/8} \end{aligned}$$

where $\mathcal{B} = \mathcal{A} \times \{\{1\}, \dots, \{M\}\}$ and $s(\mathcal{B}, n)$ is the maximum over all $(x_1, y_1), \dots, (x_n, y_n)$ in $\mathbb{R}^m \times \{1, \dots, M\}$ of the number of different sets in $\{(x_1, y_1) \cup \dots \cup (x_n, y_n) \cap B | B \in \mathcal{B}\}$. If \mathcal{A} is the class of all linear halfspaces of \mathbb{R}^m and $M=1$, then $s(\mathcal{B}, n) < 1+n^m$ by a theorem of Cover⁸ (see also Vapnik and Chervonenkis⁹). It is clear that if \mathcal{A} is the class of all intersections of M^*-1 or less linear halfspaces and $M=1$, then $s(\mathcal{B}, n) \leq (1+n^m)^{M^*-1}$. If $M>1$, then $s(\mathcal{B}, n) \leq M(1+n^m)^{M^*-1}$. Indeed, if s_1 is the number of different sets in $\{(x_1, y_1) \cup \dots \cup (x_n, y_n) \cap B | B \in \mathcal{B} \times \{\{1\}, \dots, \{M\}\}$, then the number of different sets in $\{(x_1, y_1) \cup \dots \cup (x_n, y_n) \cap B | B \in \mathcal{B} \times \{\{1\}, \dots, \{M\}\}$ is at most $M s_1$. Thus we have shown that

$$P(|L_n - \hat{L}_n| \geq \epsilon) \leq 4M(1+(2n)^m)^{M-1} e^{-n\epsilon^2/8M^2}.$$

Q.E.D.

Proof of Theorem 2.

Let us use the notation of Theorem 1 where we let A_1, \dots, A_K be the partition of \mathbb{R}^m that is generated by the nearest-neighbor rule with $(Y_1, \xi_1), \dots, (Y_K, \xi_K)$ (i.e., A_i is the set on which we estimate the state of X by ξ_i and for which Y_i is the nearest neighbor to X among Y_1, \dots, Y_K), then

$$L_n = \sum_{i=1}^K v(A_i \times \{\xi_i\}^c)$$

and

$$\hat{L}_n = \sum_{i=1}^K v_n(A_i \times \{\xi_i\}^c).$$

Thus, arguing as in Theorem 1, we have

$$|L_n - \hat{L}_n| \leq K \sup_{\substack{A \in \mathcal{A} \\ 1 \leq i \leq M}} |v_n(A \times \{i\}) - v(A \times \{i\})|$$

where \mathcal{A} is the class of all sets that are intersections of (k_n-1) or less linear halfspaces of \mathbb{R}^m . Since $K \leq k_n$, we have by the argument of Theorem 1 that

$$P(|L_n - \hat{L}_n| \geq \epsilon) \leq 4M(1+(2n)^m)^{k_n-1} e^{-n\epsilon^2/8k_n^2}.$$

Q.E.D.

Proof of Theorem 3.

It is clear that

$$\begin{aligned} |L_n - \hat{L}_n| &= |v(\bigcup_{\ell=1}^{k_n} (A_\ell \times \{\xi_\ell\}^c)) - v_n(\bigcup_{\ell=1}^{k_n} (A_\ell \times \{\xi_\ell\}^c))| \\ &= |v(\bigcup_{\ell=1}^{k_n} (A_\ell \times \{\xi_\ell\})) - |v_n(\bigcup_{\ell=1}^{k_n} (A_\ell \times \{\xi_\ell\}))| \\ &\leq \sum_{\ell=1}^{k_n} \sup_{1 \leq i \leq M} |v(A_\ell \times \{i\}) - v_n(A_\ell \times \{i\})|. \end{aligned}$$

Thus, if \mathcal{C}_ℓ is the class of sets of the form $A_\ell \times \{i\}$, $1 \leq i \leq M$, then we know by an inequality of Vapnik and Chervonenkis⁹ that

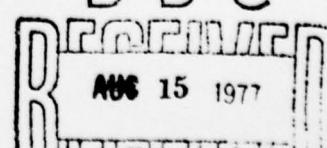
$$P(|L_n - \hat{L}_n| \geq \epsilon)$$

$$\leq k_n \sup_{1 \leq \ell \leq k} P(\sup_{C \in \mathcal{C}_\ell} |v(C) - v_n(C)| \geq \epsilon/k_n)$$

$$\leq 4k_n \left(\sup_{1 \leq \ell \leq k_n} s(\mathcal{C}_\ell, 2n) \right) e^{-n(\epsilon/k_n)^2/8}$$

$$\leq 4k_n \min(1+2n, M) e^{-n\epsilon^2/8k_n^2}.$$

Also,



Aug 15 1977

$$P(|L_n - \hat{L}_n| \geq \epsilon)$$

$$\leq P\left\{\sum_{\ell=1}^{k_n} \sum_{i=1}^M |v(A_{\ell} \times \{i\}) - v_n(A_{\ell} \times \{i\})| \geq \epsilon\right\}$$

$$\leq k_n M \sup_{\substack{1 \leq i \leq M \\ 1 \leq \ell \leq k_n}} P(|v(A_{\ell} \times \{i\}) - v_n(A_{\ell} \times \{i\})| \geq \epsilon/k_n M)$$

$$\leq 2k_n M e^{-2n\epsilon^2/M^2 k_n^2}$$

by an inequality of Hoeffding¹⁵. Furthermore,

$$|L_n - \hat{L}_n|$$

$$\leq \sup_{\text{all } (\lambda_1, \dots, \lambda_{k_n})} \left| v\left(\bigcup_{\ell=1}^{k_n} (A_{\ell} \times \{\lambda_{\ell}\})\right) - v_n\left(\bigcup_{\ell=1}^{k_n} (A_{\ell} \times \{\lambda_{\ell}\})\right) \right|$$

$$\text{from } \{1, \dots, M\}^{k_n}$$

so that

$$P(|L_n - \hat{L}_n| \geq \epsilon) \leq 4s(\mathcal{D}^*, 2n) e^{-n\epsilon^2/8}$$

where \mathcal{D}^* is the class of all sets of the form

$$\bigcup_{\ell=1}^{k_n} (A_{\ell} \times \{\lambda_{\ell}\}) \text{ where } (\lambda_1, \dots, \lambda_{k_n}) \in \mathcal{D} = \{1, \dots, M\}^{k_n}$$

Clearly, $s(\mathcal{D}^*, 2n) \leq 2^{2n}$ for all k_n . However, if

$k_n < 2n$, then $s(\mathcal{D}^*, 2n) \leq M^{k_n}$ and, in general, we must have that $s(\mathcal{D}^*, 2n) \leq 2^{k_n} (2n/k_n)^{k_n}$. This proves the inequality with g_{n3} .

Finally, notice that

$$P(|L_n - \hat{L}_n| \geq \epsilon)$$

$$\leq \sum_{\mathcal{D} \in \mathcal{D}^*} P\left\{ \left| v\left(\bigcup_{\ell=1}^{k_n} (A_{\ell} \times \{\lambda_{\ell}\})\right) - v_n\left(\bigcup_{\ell=1}^{k_n} (A_{\ell} \times \{\lambda_{\ell}\})\right) \right| \geq \epsilon/M k_n \right\}$$

$$\leq 2M e^{-2n\epsilon^2/M^2 k_n^2}$$

Q.E.D.

Proof of (1).

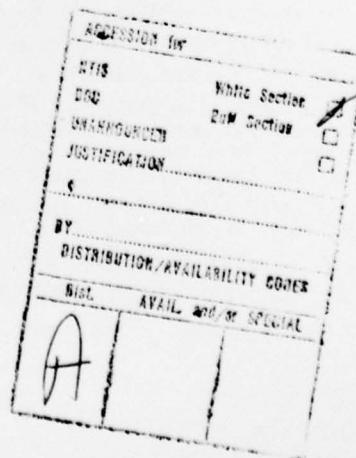
Inequality (1) is a corollary of Hoeffding's inequality¹⁵ if we note that $|L_n - \hat{L}_n| = |v(C) - v_n(C)|$ where

$$C = \bigcup_{\ell=1}^{\infty} (A_{\ell} \times \{\xi_{\ell}\})^C$$

Q.E.D.

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